

# Algebraic Optimization Degree

A Macaulay2 package

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# Introduction

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## Setting

Maximize/minimize an objective function  $\Psi$  over an algebraic variety  $X$ .

The algebraic degree of an optimization problem gives an algebraic measure of complexity of a problem

The algebraic degree of an optimization problem is an important invariant in applied algebraic geometry:

- nearest point problems (Draisma et al. 2016),
- maximum likelihood estimation (Catanese et al. 2006; Hoşten, Khetan, and Sturmfels 2005),
- semidefinite programming (Graf von Bothmer and Ranestad 2009).

## A simple example

Let  $Y$  be a binomial random variable with 2 trials.

The probability mass function of  $Y$  is given by the map

$$p \mapsto ((1 - p)^2, 2p(1 - p), p^2) =: (p_0, p_1, p_2)$$

The Zariski closure of the image is the algebraic variety corresponding to the ideal

$$I = \langle 4p_0p_2 - p_1^2, 1 - p_0 - p_1 - p_2 \rangle \subseteq \mathbb{R}[p_0, p_1, p_2]$$

## A simple example

Suppose we observe  $Y$  and collect the results in a vector  $(u_0, u_1, u_2)$ .

The likelihood function is

$$\Psi(p_0, p_1, p_2) = p_0^{u_0} p_1^{u_1} p_2^{u_2}$$

### Maximum likelihood estimation

Maximize  $\Psi$  subject to  $(p_0, p_1, p_2) \in \mathbb{V}(I)$  and  $0 \leq p_i \leq 1$  for  $i \in \{1, 2, 3\}$ .

Approach: find all complex critical points of  $\Psi$ . The number of critical points is the *maximum likelihood degree*.

## General setting

- $X \subset \mathbb{C}^n$  an affine variety of codimension  $c$
- $\Psi: X \rightarrow \mathbb{C}$  an objective function, with gradient  $\nabla\Psi$

### Definition

The **critical ideal**  $\text{Crit}_0(\Psi, X)$  is the ideal of the set of isolated critical points of  $\Psi$  on the regular locus of  $X$ . The degree of  $\text{Crit}_0(\Psi, X)$  is the **optimization degree**.

$$S := \left( \langle f_1, \dots, f_N \rangle + \left\langle (c+1) \times (c+1) \text{ minors of } \begin{bmatrix} \nabla\Psi \\ \nabla f_1 \\ \vdots \\ \nabla f_N \end{bmatrix} \right\rangle \right) : I_{X_{\text{sing}}}^\infty,$$

where  $f_1, \dots, f_n$  generate the radical ideal of  $X$ . Let  $P$  denote the ideal of positive dimensional components of the variety of  $S$ . Then,

$$\text{Crit}_0(\Psi, X) = S : P^\infty.$$

## Euclidean distance degree

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# ED-degree

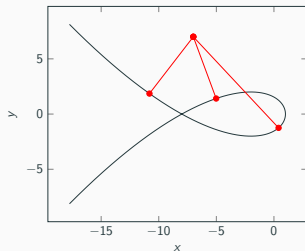
## Problem

Given a variety  $X \subseteq \mathbb{C}^n$  and a point  $u \in \mathbb{C}^n$ , find the point  $x \in X$  that minimizes the Euclidean distance  $\|x - u\|_2$ .

The number of critical points of the Euclidean distance function for a *generic* point  $u$  is the **Euclidean distance degree**

## Example

```
i1 : R = QQ[x,y];  
i2 : I = ideal(27*y^2-(1-x)*(8+x)^2);  
i3 : probabilisticEDDegree(I)  
o3 = 5  
i4 : symbolicEDDegree(I)  
o4 = 5
```





## ED-degree via projections

Let  $X \subseteq \mathbb{P}^n$  have codimension  $\geq 2$ . Let  $\pi: \mathbb{P}^n \rightarrow \mathbb{P}^{n-1}$  be a rational map defined by a general linear map  $\mathbb{C}^{n+1} \rightarrow \mathbb{C}^n$ .

### Theorem (Draisma et al. (2016))

*The ED-degree of  $X$  is equal to the ED-degree of  $\pi(X)$ .*

### Example

```
i1 : R = ZZ/101[x_1..x_6]
i2 : M = genericSymmetricMatrix(R,3)
i3 : I = minors(2,M)
i4 : elapsedTime probabilisticEDDegree I
    -- 566.658 seconds elapsed
o4 = 13
i5 : elapsedTime projectionEDDegree I
    -- 0.115337 seconds elapsed
o5 = 13
```

## Maximum likelihood degree

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# Introduction

- Probability simplex

$$\Delta_n := \{(p_0, \dots, p_n) : p_i \geq 0, \sum_i p_i = 1\}$$

- The statistical model is the solution set of homogeneous polynomials in  $p_0, \dots, p_n$ .
- Let  $V$  be the variety corresponding to the model.

## Maximum Likelihood Estimation

Given a vector  $u \in \mathbb{N}^{n+1}$  of observations, maximize

$$L = \frac{p_0^{u_0} + \dots + p_n^{u_n}}{(p_0 + \dots + p_n)^{u_0 + \dots + u_n}}$$

subject to  $(p_0, \dots, p_n) \in V \cap \Delta_n$ .

## Definition (Hoşten, Khetan, and Sturmfels (2005))

The **ML-degree** is the number of critical points of  $L$  for generic  $u$ .

# Implicit form

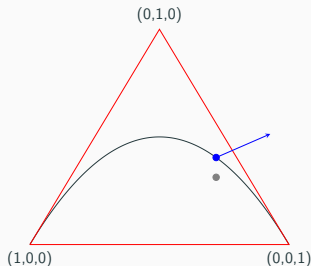
## Example

Binomial random variable with 2 trials. Say we observe  $u = (2, 5, 9)$ .

```
i1 : R = QQ[p_0..p_2]
i2 : P = ideal(4*p_0*p_2-p_1^2)
i3 : MLequationsIdeal(P, {2,5,9})

o3 = ideal (23p1 - 18p2 , 529p0 - 81p2 )

i4 : MLequationsDegree(P)
o4 = 1
```



## Definition

The **scaled toric variety** is the Zariski closure of the map  $\phi_{A,c} : (\mathbb{C}^*)^d \rightarrow (\mathbb{C}^*)^r$  in  $\mathbb{C}^r$  given by

$$\phi_{A,c}(\theta_1, \dots, \theta_d) = (c_1\theta^{a_1}, \dots, c_r\theta^{a_r}).$$

## Example

```
i1 : A = matrix {{1,1,1,0,0,0,0,0,0}, {0,0,0,1,1,1,0,0,0},
{0,0,0,0,0,0,1,1,1},{1,0,0,1,0,0,1,0,0},{0,1,0,0,1,0,0,1,0}}
i2 : c = {1,2,3,1,1,1,1,1,1};
i3 : elapsedTime toricMLDegree(A, c)
-- 0.00849565 seconds elapsed
o3 = 3
```

The corresponding computation using `MLequationsDegree` took 228 seconds.

## Optimizations and future

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# Fritz John and Lagrange multipliers

- Let  $A \in \mathbb{C}[x]^{m \times n}$  be a (generically) full-rank matrix,  $m \geq n$ .
- Computing  $\langle n \times n \text{ minors of } A \rangle$  may be difficult.
- Lagrange/Fritz John:

$$\begin{bmatrix} z_1 & \cdots & z_n \end{bmatrix} A = 0$$
$$\sum_i c_i z_i = 1$$

## Example

```
i1 : R = QQ[x_1..x_6];
i2 : M = genericSymmetricMatrix(R,3);
i3 : I = minors(2, M);
i4 : probabilisticEDDegree I
    -- 179.771 seconds elapsed
o4 = 13
i5 : probabilisticFritzJohnEDDegree I
    -- 4.95628 seconds elapsed
o5 = 13
```

Future: push beyond the boundaries of symbolic methods with numerical algebraic geometry

- Compute all critical points
- Find the optimal point
- Parametric homotopy
- MonodromySolver



# Conclusion

## Package features

- General optimization problem
  - Ideal
  - Lagrange
- Euclidean Distance Degree
  - Probabilistic
  - Symbolic
  - Projective
    - Multidegree
    - Projections
    - Sections
  - Fritz John
- Maximum Likelihood degree
  - Symbolic
  - Parametric
  - Toric

Package available at  
[https://github.com/Macaulay2/  
Workshop-2020-Cleveland/tree/  
ISSAC-AlgOpt/alg-stat/  
AlgebraicOptimization](https://github.com/Macaulay2/Workshop-2020-Cleveland/tree/ISSAC-AlgOpt/alg-stat/AlgebraicOptimization)

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**For more questions, join the link**  
**<https://bluejeans.com/550710311>**  
**on Jul 21, 18:00 EET**

## References




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Catanese, Fabrizio et al. (2006). “The maximum likelihood degree”. In: *Amer. J. Math.* 128.3, pp. 671–697. ISSN: 0002-9327. URL: [http://muse.jhu.edu/journals/american\\_journal\\_of\\_mathematics/v128/128.3catanese.pdf](http://muse.jhu.edu/journals/american_journal_of_mathematics/v128/128.3catanese.pdf).



Draisma, Jan et al. (2016). “The Euclidean distance degree of an algebraic variety”. In: *Found. Comput. Math.* 16.1, pp. 99–149. ISSN: 1615-3375. DOI: 10.1007/s10208-014-9240-x. URL: <https://doi.org/10.1007/s10208-014-9240-x>.

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-  Holme, Audun (1988). “The geometric and numerical properties of duality in projective algebraic geometry”. In: *Manuscripta Math.* 61.2, pp. 145–162. ISSN: 0025-2611. DOI: 10.1007/BF01259325. URL: <https://doi.org/10.1007/BF01259325>.
-  Hoşten, Serkan, Amit Khetan, and Bernd Sturmfels (2005). “Solving the likelihood equations”. In: *Found. Comput. Math.* 5.4, pp. 389–407. ISSN: 1615-3375. DOI: 10.1007/s10208-004-0156-8. URL: <http://dx.doi.org/10.1007/s10208-004-0156-8>.

**Backup slides**

# ED-degree via multidegree

## Theorem (Draisma et al. (2016))

*The ED-degree of  $X$  is the sum of the polar classes of  $\mathcal{N}$*

### Example

```
i1: R = QQ[x_0..x_3]
i2: J = ideal det(matrix{
    {x_0, x_1, x_2},
    {x_1, x_0, x_3},
    {x_2, x_3, x_0}});
i3: symbolicMultidegreeEDDegree J
    -- 0.263975 seconds elapsed
o3 = 13
i4: probabilisticMultidegreeEDDegree J
    -- 0.264061 seconds elapsed
o4 = 13
```

# Projective varieties

The ED-degree of a projective variety  $X \subseteq \mathbb{P}^n$  is the ED-degree of its affine cone in  $\mathbb{C}^{n+1}$ .

## Definition

- The **conormal variety**  $\mathcal{N}$  is the closure of the set

$$\{(x, u) \in \mathbb{P}^n \times (\mathbb{P}^n)^* : x \in X \setminus X_{\text{sing}}, u \in T_x X\}$$

- We say  $X$  is in **general position** if  $\mathcal{N}$  does not intersect the diagonal  $\Delta \subseteq \mathbb{P}^n \times (\mathbb{P}^n)^*$ .

Throughout the rest of the section,  $X$  is a projective variety in general position.

# ED-degree via linear sections

## Theorem

$$\text{EDdegree}(X) = \begin{cases} \text{EDdegree}(X \cap H), & \text{if } \text{codim}(X^*) \geq 2 \\ \text{EDdegree}(X \cap H) + \deg(X^*), & \text{if } \text{codim}(X^*) = 1 \end{cases}$$

## Notes

- Requires computing the **projective dual**  $X^*$ , i.e. the projection of  $\mathcal{N}$  onto its second factor.
- Iterated application requires computation of  $(X \cap H)^*$

## Theorem (Holme (1988))

*Let  $H$  be a hyperplane. If  $X^*$  has codimension  $\geq 2$ , then*

$$(X \cap H)^* = \pi_H(X^*),$$

*where  $\pi_H$  is the projection with center point  $H \in (\mathbb{P}^n)^*$*



# ML-degree of parametric models

In many applications, the statistical model is given as the image of a polynomial map

## Example

```
i1 : R = QQ[p]
i2 : param = {(1-p)^2, 2*p*(1-p), p^2}
i3 : parametricMLIdeal(param, {2,5,9})
o3 = ideal(32p - 23)
i4 : parametricMLDegree(param)
o4 = 1
```